Input example - File class

```
class File {
```

```
int linesInFile; int linesRead;
boolean closed; boolean lineInBuffer; boolean eof;
```

```
//@invariant linesRead >= 0 && linesRead <= linesInFile;
//@initial linesRead == 0 && linesInFile == 5
    && !closed && !lineInBuffer && !eof;
void File() { ... }
```

Input example - *File* class

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- Determining a state consists on determining the set of methods which the precondition is implied by the constructor's initial condition

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- Determining the next states is similar to determining the initial state, only it uses the postcondition of *m* instead of the initial condition

- If m is of boolean type and its postcondition specifies two states, two states are determined: One for the true result and other for the false result
- For example, the postcondition of method *eof* implies that:
 - If it returns *true*, there is no more lines to read. This state is valid for the *read* method but not for the *close* method.
 - If it returns *false*, there is at least one more line to read. This state is valid for the *close* method but not for the *read* method.

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- Its execution causes the typestate to transit into a decision state which will have two transitions, one for each possible result of *eof*:

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$$\delta(\{eof\}, eof) = \{eof_choice\}$$

$$\delta(\{eof_choice\}, true) = \{close\}$$

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Algorithm steps

States {read} and {close}, since they have not been explored yet, are added to W.

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After fully exploring the state, the algorithm then explores the state in the head of W and repeats the process until W is empty

Output example - Typestate of the File class



Algorithm steps: state id assignment

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Algorithm steps: state shared status

- The algorithm also determines the shared status of each state
- A state is considered shared if it only transits to itself or to an equivalent state

Algorithm steps: state translation

 Using the previous set and the transition relation of the typestate, each state is then translated into an usage state

 $\{(0, \{eof\}), (1, \{close\})(2, \{read\})\}$

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usage lin { File ; 0 } where
            0 = lin { eof ; <1 + 2>}
            1 = _{close; _}
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    0 = lin { eof ; <1 + 2>}
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    0 = lin { eof ; <1 + 2>}
    1 = lin { close ; end}
    2 = lin { read ; 0 }
```

Output example - Usage of the File class

```
usage lin { File ; Q1 } where
    0 = lin { eof ; <1 + 2>}
    1 = lin { close ; end}
    2 = lin { read ; 0}
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Algorithm steps - overview

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 - Analyses the code of the method and:
 - For every initialization, sets the usage state of the initialized variable with the usage state of the value
 - For every call, changes the current usage state of the object the method was called

Algorithm steps - determining the usage state of an object

When determining the usage state of an object, the algorithm checks the first usage state that has a set of method whose preconditions are implied by the assertions that specifies its state