Input example - *File* class

. . .

```
class File {
```

```
int linesInFile; int linesRead;
boolean closed; boolean lineInBuffer; boolean eof;
```

```
\frac{1}{2} //@invariant linesRead >= 0 && linesRead <= linesInFile:
//@initial linesRead == 0 && linesInFile == 5
           && !closed && !lineInBuffer && !eof;
void File () { ... }
```
Input example - *File* class

. . .

. . .

```
class File {
```

```
//@requires linesRead < linesInFile && !closed && lineInBuffer
           && !eof;
\frac{1}{2} //@ensures linesRead + 1 <= linesInFile
           && !closed && !lineInBuffer && !eof;
string read () \{ \ldots \}
```
Input example - *File* class

. . .

. . .

```
class File {
```

```
//@requires linesRead <= linesInFile && !closed && !lineInBuffer
          && !eof;
//@ensures (linesRead == linesInFile -> !lineInBuffer && eof)
          && !closed;
boolean eof() { ... }
```
Input example - *File* class

. . .

```
class File {
```

```
\frac{1}{2} //@requires linesRead == linesInFile && eof && !closed;
//@ensures linesRead == linesInFile && eof && closed;
void close() {
          close = true}
```
Algorithm steps: initialisation

■ Starts by determining the initial state.

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- Determining a state consists on determining the set of methods which the precondition is implied by the constructor's initial condition

Algorithm steps: initialisation

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Algorithm steps: iteration

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 $W = \{\{eof\}\}\$

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■ Before exploring it, the state A will be added to the set of states *S*, which is the set of states of the typestate $S = \{ \{ \text{eof} \} \}$

Algorithm steps

For each method m in state A, the algorithm will determine the state the typestate transits to when *m* is executed.

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- For each method *m* in state *A*, the algorithm will determine the state the typestate transits to when *m* is executed.
- Determining the next states is similar to determining the initial state, only it uses the postcondition of *m* instead of the initial condition

Algorithm steps

- If *m* is of boolean type and its postcondition specifies two states, two states are determined: One for the true result and other for the false result
- **For example, the postcondition of method** *eof* **implies that:**
	- If it returns *true*, there is no more lines to read. This state is valid for the *read* method but not for the *close* method.
	- If it returns *false*, there is at least one more line to read. This state is valid for the *close* method but not for the *read* method.

Algorithm steps

- **This means that after** *eof* there will be two possible states to transit to depending of the returned value
- \blacksquare Its execution causes the typestate to transit into a decision state which will have two transitions, one for each possible result of *eof*:

Algorithm steps

- **This means that after** *eof* there will be two possible states to transit to depending of the returned value
- \blacksquare Its execution causes the typestate to transit into a decision state which will have two transitions, one for each possible result of *eof*:

 δ ({*eof*}, *eof*) = {*eof choice*} δ {{*eof_choice*}, *true*} = {*close*} δ {{*eof choice*}, *false*} = {*read*}

Algorithm steps

States {*read*} and {*close*}, since they have not been explored yet, are added to *W*.

Algorithm steps

■ The algorithm does the same to every method in *A*

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After fully exploring the state, the algorithm then explores the state in the head of *W* and repeats the process until *W* is empty

Output example - Typestate of the *File* class

Algorithm steps: state id assignment

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For the *File* class this set will be: {(0, {*eof* }),(1, {*close*})(2, {*read*})}

Algorithm steps: state shared status

- \blacksquare The algorithm also determines the shared status of each state
- A state is considered shared if it only transits to itself or to an equivalent state

Algorithm steps: state translation

Using the previous set and the transition relation of the typestate, each state is then translated into an usage state

{(0, {*eof* }),(1, {*close*})(2, {*read*})}

```
\delta({\mathcal{B}}) = {\mathcal{B}} eof choice
\delta({eof choice}, true) = {close}
\delta({\mathcal{B} of\_choice}, \mathcal{B} is a) = {\mathcal{B} of}\\delta({\mathcal{A}}close),close) = {\mathcal{A}}\delta({\text{read}}, read) = {eof}
```
usage lin { File : 0} where

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```
usage lin { File ; 0} where $0 = \{ eof : \}$

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{(0, {*eof* }),(1, {*close*})(2, {*read*})}

 $\delta({\mathcal{B} of \}, \mathsf{eof}) = {\mathsf{eof_choice}}$ δ {{*eof choice*}, *true*} = {*close*} $\delta({\mathcal{B}(\mathsf{eof_choice}}), \mathsf{false}) = {\mathsf{read}}$ $\delta({\mathcal{A}}$ *close* $),$ *close* $) = {\mathcal{A}}$ $\delta({\text{read}}$, *read*) = ${\text{def}}$

usage lin { File : 0} where $0 = \{ \text{eof} \, ; \, < + \, > \}$

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{(0, {*eof* }), (**1**, {*close*}), (**2**, {*read*})}

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\delta({eof choice}, true) = {close}
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```
usage lin { File : 0} where $0 = \text{lin} \{ \text{eof} : \langle 1 + 2 \rangle \}$

Algorithm steps: state translation

Using the previous set and the transition relation of the typestate, each state is then translated into an usage state

{(0, {*eof* }),(1, {*close*}),(2, {*read*})}

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```
usage lin { File : 0} where $0 = \text{lin} \{ \text{eof} : \langle 1 + 2 \rangle \}$ $1 = \{ close; \}$

Algorithm steps: state translation

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{(0, {*eof* }),(1, {*close*}),(2, {*read*})}

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\delta({\mathcal{B}}) = {\mathcal{B}} eof choice
\delta({eof choice}, true) = {close}
\delta({\mathcal{B}(\mathsf{eof\_choice}}), \mathsf{false}) = {\mathsf{read}}\delta({close}, close) = {}
\delta({\text{read}}, read) = {eof}
```

```
usage lin { File : 0} where
               0 = \text{lin} \{ \text{eof} : \langle 1 + 2 \rangle \}1 = \text{lin} \{ \text{close} : \text{end} \}
```
Algorithm steps: state translation

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{(0, {*eof* }),(1, {*close*}),(2, {*read*})}

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Algorithm steps: state translation

Using the previous set and the transition relation of the typestate, each state is then translated into an usage state

{(**0**, {*eof*}),(1, {*close*}),(2, {*read*})}

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```
usage lin { File : 0} where
                0 = \text{lin} \{ \text{eof} : \langle 1 + 2 \rangle \}1 = \text{lin} \{ \text{close} : \text{end} \}2 = \text{lin} \{ \text{read} : 0 \}
```
Output example - Usage of the *File* class

```
usage lin { File : Q1 } where
        0 = \text{lin} \{ \text{eof} : \langle 1 + 2 \rangle \}1 = \text{lin} \{ \text{close} : \text{end} \}2 = \text{lin} \{ \text{read} : 0 \}
```
Algorithm steps - overview

Goes through all the methods of each class (following the order of the usage) and, for each one:

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	- **1** Determines the usage state of every parameter using the precondition of the method

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	- **1** Determines the usage state of every parameter using the precondition of the method
	- 2 Determines the usage state of the return type using the postcondition of the method

- Goes through all the methods of each class (following the order of the usage) and, for each one:
	- **1** Determines the usage state of every parameter using the precondition of the method
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	- **3** Analyses the code of the method and:

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		- For every initialization, sets the usage state of the initialized variable with the usage state of the value

- Goes through all the methods of each class (following the order of the usage) and, for each one:
	- **1** Determines the usage state of every parameter using the precondition of the method
	- 2 Determines the usage state of the return type using the postcondition of the method
	- **3** Analyses the code of the method and:
		- For every initialization, sets the usage state of the initialized variable with the usage state of the value
		- For every call, changes the current usage state of the object the method was called

Algorithm steps - determining the usage state of an object

When determining the usage state of an object, the algorithm checks the first usage state that has a set of method whose preconditions are implied by the assertions that specifies its state